

RSA®Conference2021

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RESILIENCE

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SoK: How (not) to Design and Implement Post-Quantum Cryptography

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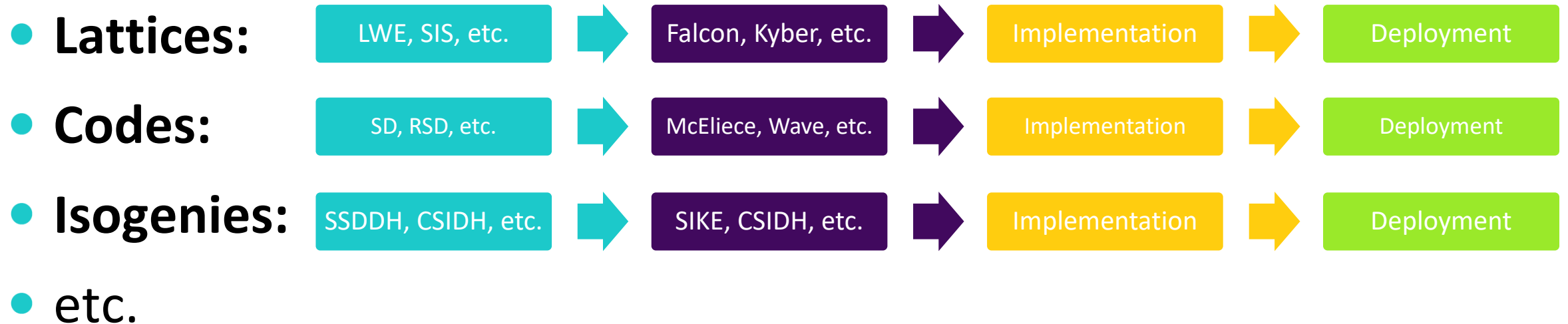
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This SoK: a transversal survey of post-quantum cryptography

Most works focus on (one aspect of) one family:



We tried to abstract away the family, and focus on the process:



Plan

- **This talk:**



- **The paper:**



- **The ePrint:**



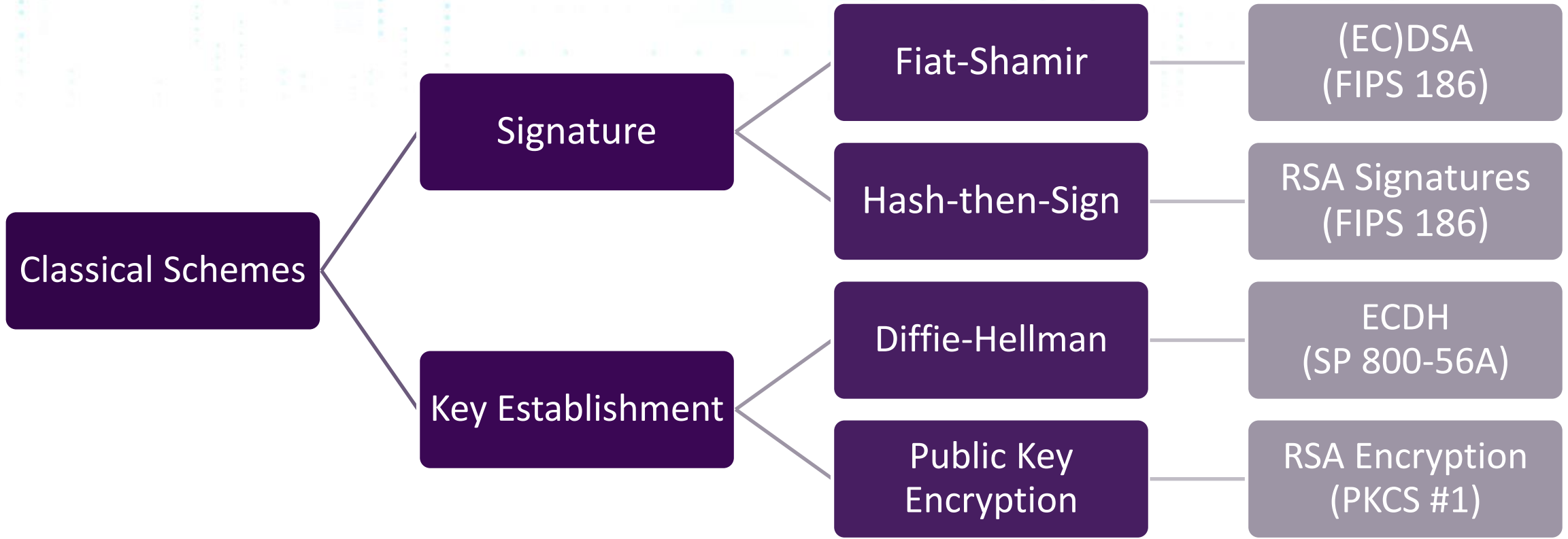
Our goals

- Survey essential works
- Establish trends and patterns
- Provide “lessons learned”

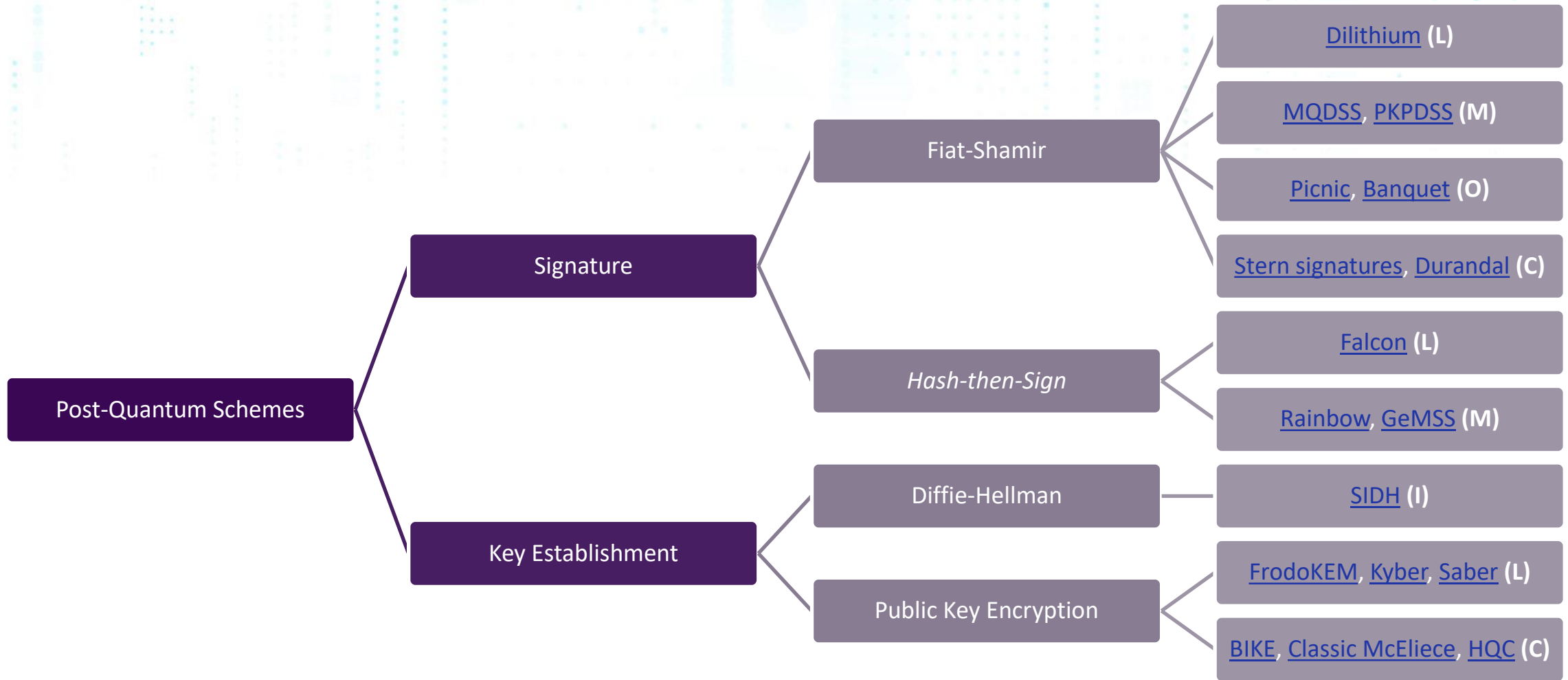


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Theory and Design of Schemes



CLASSICAL SCHEMES



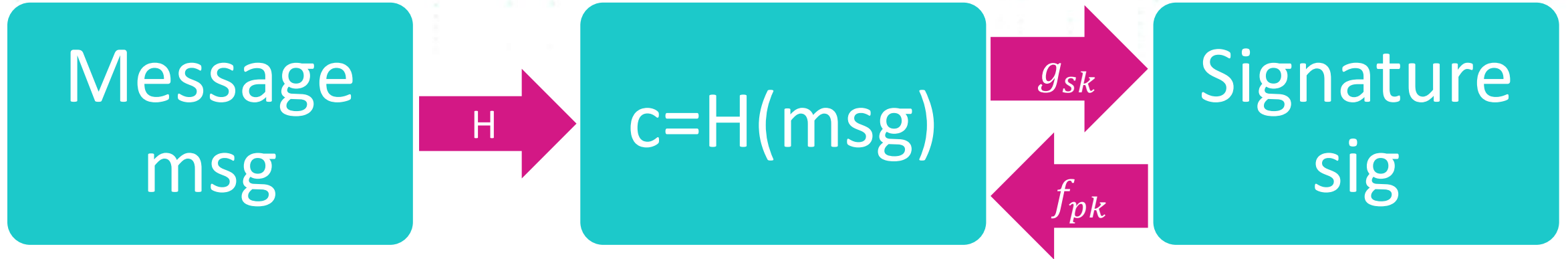
**MOST POST-QUANTUM SCHEMES -
(L)ATTICES, (C)ODES, (I)SOGENIES, (M)ULTIVARIATE, (O)NE-WAY FUNCTIONS**



“Just apply one of these four paradigms to your favourite problem”? Not so simple.

- These paradigms are useful guidelines, but they are **no panacea**.
 - Rigidly applying a paradigm may result in an **inefficient scheme** or a **broken assumption**.
 - It is more important to preserve the **security proof** than the paradigm.
- Most efficient schemes tweak paradigms to fit the assumption.
 - Full-domain hash **without permutations** (see [next slide](#))
 - Fiat-Shamir **with aborts** and further refinements [[Lyu09](#),[BG14](#),[Dilithium](#)]
 - Various soundness-amplification tricks [[DG19](#),[KKW18](#)]See the paper for a complete discussion.

Example: Full-Domain Hash Signatures



We say that the pair $(f_{pk}: X \rightarrow Y, g_{sk}: Y \rightarrow X)$ is a trapdoor permutation if:

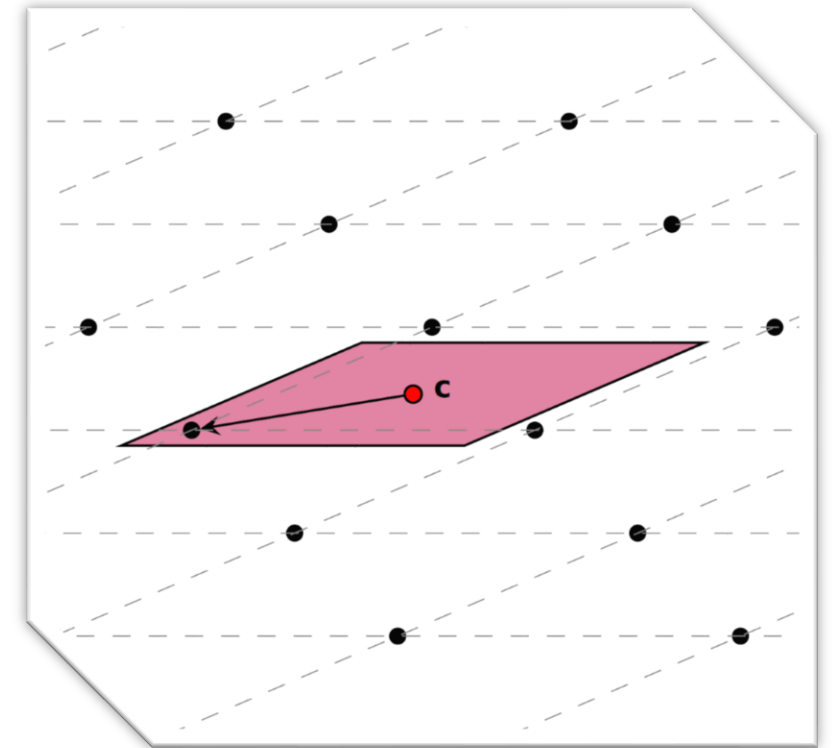
1. Given only pk , it is computationally hard to invert f_{pk} on (almost) all inputs.
2. $f_{pk} \circ g_{sk}$ is the identity over Y , and $X = Y$ (hence f_{pk} and g_{sk} are permutations).

Canonical example: RSA signatures [[RSA78](#)] and its many variants.

Provable security is well-studied [[BR96](#), [Coron00](#)].

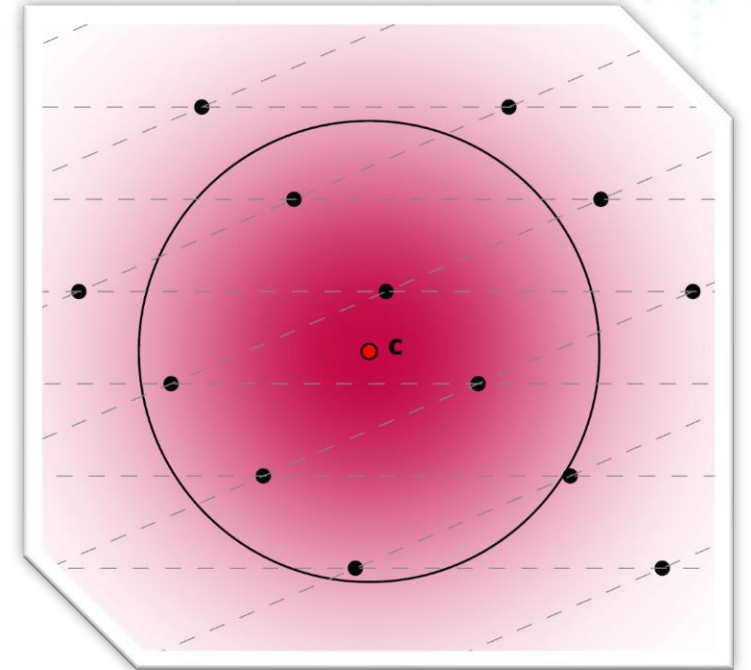
Can we transpose this to the post-quantum setting?

- Initial attempts: GGHSign [[GGH97](#)](lattices), [[CFS01](#)](codes), NTRUSign [[HHPSW03](#)] (lattices)
 - **CFS'01**: poor scalability of parameters
 - **GGHSign and NTRUSign**:
 - (f_{pk}, g_{sk}) was not a trapdoor permutation.
 - The usual security proof did no longer apply
 - Worse, each signature leaked information about the signing key sk , leading to practical attacks [[NR06](#), [DN12](#)].



Solution: relaxing the notion of trapdoor permutation

- Gentry-Peikert-Vaikuntanathan [[GPV08](#)]:
 - Trapdoor preimage sampleable functions (TPSF)
 - Weaker than trapdoor permutation
 - Can be instantiated from lattices (but not codes)
 - Still strong enough for a security proof
- A further relaxation [[DST19](#),[CGM19](#),[CD20](#)]:
 - Average TPSF
 - Can be instantiated from codes and lattices
- Trapdoor permutation \Rightarrow TPSF \Rightarrow Average TPSF ($\Rightarrow ?$)
- Examples: Falcon [[FHK⁺17](#)] (TPSF, lattices), Wave [[DST19](#)] (Average TPSF, codes)



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Implementation

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An Overview of PQC Implementations

- Transitioning to PQC will be tough.
- PQC will bring many new and unique challenges.
 - Larger public/secret keys, signs, ciphertexts.
 - Needs more resources; time, hardware, energy.
 - New operations and paradigms.
- Implementations will be more complex.
 - Rejection sampling,
 - Sampling non-uniform distributions,
 - Decryption failures, etc.



Implementation Attacks on PQC

- Many initial implementations not isochronous.
 - Timing leakage [[Str10](#),[AHP⁺12](#),[Str13](#),[ELP⁺18](#),[EFG⁺17](#)].
 - FLUSH+RELOAD cache attacks [[BHL⁺16](#),[PBY17](#)].
 - Data-dependent branching / branch tracing [[EFG⁺17](#)].
- Recent attacks exploit implementation mistakes.
 - Non-isochronous memcmp() in FO transform [[GJN20](#)].
 - Errors in domain separation in FO [[BDG20](#)].
- Fixable by following secure coding practices.



Side-Channel Attacks on PQC

- A talk on this topic has been done [[AH21](#)].
- Power analysis targets secrets via:
 - Matrix multiplication [[ATT+18](#),[PSK+18](#),[BFM+19](#)].
 - Polynomial multiplication [[HCY20](#)].
 - Syndrome decoding [[RHH+17](#),[SKC+19](#)].
- Also, fault, cold-boot, and key reuse attacks.
- Side-channel “hints” for security evaluations.
 - For lattice reduction [[DDG+20](#)] and ISD [[HPR+21](#)].



Masking and Hiding in PQC

- We've only just begun protecting these schemes.
- So far, only masked Saber, Kyber and Dilithium.
 - First-order [[MGT+19](#)], higher-order [[BDK+19](#),[BGR+21](#)].
- Countermeasures also not a guarantee.
 - QcBits masking [[RHH+17](#)] was attacked [[SKC+19](#)].
 - Masked comparison [[BPO+20](#)] was attacked [[BDH+21](#)].
- Hedging mitigates faults for Fiat-Shamir signs [[AOT+20](#)].



Benchmarking PQC

- Evaluate performances in SW, HW, and SCA.
 - On ARM Cortex M4 and Xilinx Artex 7 FPGA.
- Plenty of repos exist for SW benchmarking.
 - [PQClean](#), [pqm4](#), [SUPERCOP](#), etc.
- Seed expanding in pqm4 can take >50% runtime [[KRS⁺19](#)].
- Lots of HW designs exist, too.
 - With a large variety in resources/performance.



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Thank you!

Full paper: <https://eprint.iacr.org/2021/462>